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Robust Power System State Estimation From Rank-One Measurements

Gang Wang[®], *Member, IEEE*, Hao Zhu, *Member, IEEE*, Georgios B. Giannakis[®], *Fellow, IEEE*, and Jian Sun[®], *Member, IEEE*

Abstract—The unique features of current and upcoming 5 energy systems, namely high penetration of uncertain re-6 newables, unpredictable customer participation, and pur-7 poseful manipulation of meter readings, all highlight the 8 9 need for fast and robust power system state estimation (PSSE). In the absence of noise, PSSE is equivalent to 10 solving a system of quadratic equations, which, also re-11 lated to power flow analysis, is NP-hard in general. Assum-12 ing the availability of all power flow and voltage magnitude 13 measurements, this paper first suggests a simple algebraic 14 technique to transform the power flows into rank-one mea-15 surements, for which the ℓ_1 -based misfit is minimized. To 16 uniquely cope with the nonconvexity and nonsmoothness 17 18 of ℓ_1 -based PSSE, a deterministic proximal-linear solver is developed based on composite optimization, whose gener-19 alization using stochastic gradients is discussed too. This 20 paper also develops conditions on the ℓ_1 -based loss func-21 tion such that exact recovery and quadratic convergence of 22 23 the proposed scheme are guaranteed. Simulated tests using several IEEE benchmark test systems under different set-24 25 tings corroborate our theoretical findings, as well as the fast convergence and robustness of the proposed approaches. 26

Index Terms—Bad data analysis, composite optimization,
 least-absolute-value (LAV) estimator, proximal-linear algo rithm, supervisory control and data acquisition (SCADA)
 measurement.

I. INTRODUCTION

HE North American power grid is praised as the greatest engineering achievement of the 20th century [1]. To maintain grid efficiency, reliability, and sustainability, system

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G. Wang and G. B. Giannakis are with the Digital Technology Center and the Department of Electrical and Computer Engineering, University of Minnesota, Minneapolis, MN 55455 USA (e-mail: gangwang@ umn.edu; georgios@umn.edu).

H. Zhu is with the Department of Electrical and Computer Engineering, The University of Texas at Austin, Austin, TX 78712 USA (e-mail: haozhu@utexas.edu).

J. Sun is with the State Key Laboratory of Intelligent Control and Decision of Complex Systems and the School of Automation, Beijing Institute of Technology, Beijing 100081, China (e-mail: sunjian@bit.edu.cn).

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operators constantly monitor the operating conditions of electricity networks [2], [3]. In the 1960s, power system engineers tried to compute voltages at critical buses based on meter readings manually collected from geometrically distributed current and potential transformers. Due to timing, model mismatches, and metering errors, however, the exact (ac) power flow equations were never infeasible.

With the development of supervisory control and data acqui-42 sition (SCADA) systems, a wealth of improved data metered 43 from across the network became available. In the seminal work 44 of Schweppe et al. [4], the modern statistical foundation for 45 power system state estimation (PSSE) was laid. Given a col-46 lection of SCADA data along with corresponding measurement 47 matrices, the goal of PSSE is to compute the complex voltages 48 (or, the voltage magnitudes and angles if polar coordinates are 49 used) at all network buses. Since then, substantial contributions 50 have been devoted to PSSE. Interested readers can refer to [3] 51 for a review of recent developments on PSSE. 52

Based on the weighted least-squares (WLS) estimation cri-53 terion, the Gauss-Newton solver is arguably the "workhorse" 54 for PSSE, and it is also employed in practice [2]. Yet, the non-55 convex nature of WLS poses challenges on the Gauss-Newton 56 method, including sensitivity to initialization and outliers, as 57 well as no convergence guarantee [5]. To address these chal-58 lenges, semidefinite programming (SDP) relaxation approaches 59 have been pursued [6]–[8]. However, SDP incurs computational 60 complexity that does not scale well with problem dimension, 61 discouraging its use in practical settings. 62

With utilities increasingly shifting toward smart grid tech-63 nology and other upgrades with inherent cyber vulnerabilities, 64 correlative threats from adversarial cyberattacks on the North 65 American power grid continue to grow in frequency and form 66 [9]. These introduce new yet critical challenges to PSSE, partic-67 ularly to the WLS-based SE solvers, concerning data integrity 68 and uninformed model changes [10]–[14]. Such concerns moti-69 vate well the development of accurate and robust approaches to 70 endow PSSE with resilience to anomalous (i.e., bad) data and 71 model inaccuracies. 72

A. Related Work

In this context, robust PSSE has recently received renewed 74 interest. To cope with the malicious data, the largest normalized 75 residual (LNR) test was incorporated while performing PSSE 76 [10]. The least-median-squares and least-trimmed-squares 77

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based alternatives were pursued [15]. Unfortunately, the aforementioned robust PSSE proposals incur unfavorable computational complexities and/or stringent storage requirements, which

limit their practical uses in real-world power networks.

On the other hand, the ℓ_1 -based criterion has been well known 82 in optimization and statistics for its robustness to outliers [16], 83 [2, Ch. 6]. In addition to being robust, the ℓ_1 -based estimator is 84 also statistically optimal in the maximum likelihood sense, when 85 the independent additive noise follows a Laplacian distribution. 86 87 In the PSSE literature, the ℓ_1 -based criterion was advocated for bad data identification and rejection in [17]. Research focus has 88 shifted toward devising efficient and user-friendly algorithms to 89 handle the nonconvexity and nonsmoothness issues of ℓ_1 -loss 90 function; see, e.g., [18] and [19]. 91

92 B. This Paper

This paper revisits the ℓ_1 -based robust PSSE with a focus on 93 development of efficient algorithms and theory on exact state 94 recovery in the noiseless case. Leveraging recent advances in 95 96 solving rank-one quadratic equations (i.e., phase retrieval) [20], [21], we first suggest a simple algebraic procedure to transform 97 the power flows into rank-one measurements, namely with cor-98 responding transformed measurement matrices being rank one. 99 Subsequently, we develop two efficient and easy-to-implement 100 algorithms to optimize the ℓ_1 -loss of the obtained rank-one mea-101 surements. With appropriate conditions on the ℓ_1 -loss function, 102 we establish exact recovery as well as quadratic convergence 103 for our approach. Simulated tests using three IEEE benchmark 104 systems showcase the robustness and computational efficiency 105 of our proposed scheme relative to competing Gauss-Newton 106 method. 107

The rest of this paper is organized as follows. System modlog eling and problem formulation are given in Section II. The procedure to obtain rank-one measurements is presented in Section III, followed by two algorithms in Section IV. Exact state recovery and convergence are established in Section V. Numerical tests are provided in Section VI, and this paper is concluded in Section VII.

115 Notation: Matrices (column vectors) are denoted by upper-116 (lower-) case boldface letters; e.g., **A** (a). Sets are denoted 117 using calligraphic letters. Symbol $^{\mathcal{T}}(^{\mathcal{H}})$ represents (Hermitian) 118 transpose, and $\overline{(\cdot)}$ complex conjugate, whereas $\Re(\cdot)(\Im(\cdot))$ takes 119 the real (imaginary) part of a complex number.

120 II. SYSTEM MODELING AND PROBLEM FORMULATION

121 A. System Modeling

122 Consider an electric power grid modeled as a graph $\mathcal{G} =$ 123 $(\mathcal{N}, \mathcal{L})$, whose nodes $\mathcal{N} := \{1, 2, ..., N\}$ correspond to buses 124 and whose edges $\mathcal{L} := \{(n, n')\} \subseteq \mathcal{N} \times \mathcal{N}$ correspond to lines. 125 Throughout this paper, all analysis pertains to the per unit 126 (p.u.) system. The complex voltage per bus $n \in \mathcal{N}$ can be 127 given in rectangular coordinates as $v_n = \Re(v_n) + j \Im(v_n)$. For 128 brevity, all nodal voltages are stacked up to form the vector 129 $\mathbf{v} := [v_1 \cdots v_N]^T \in \mathbb{C}^N$. In the ac-based SE literature, a subset of following system variables can be measured by SCADA 130 [2, Ch. 2]: 131

- 1) $|v_n|$: the voltage magnitude at bus n; 132
- 2) P_{nn'} (Q_{nn'}): the active (reactive) power flow from buses 133 n to n' at the sending terminal;
 134
- 3) $P_n(Q_n)$: the active (reactive) power injection into bus n. 135

Compliant with the ac power flow model [2], these system 136 variables can be expressed as quadratic functions of **v**. This 137 justifies why the voltage vector **v** is referred to as the system 138 state. To this end, observe that the squared voltage magnitudes 139 $V_n := |v_n|^2 = [\Re(v_n)]^2 + [\Im(v_n)]^2$ can be written as 140

$$V_n = \mathbf{v}^{\mathcal{H}} \mathbf{H}_n^V \mathbf{v}, \text{ with } \mathbf{H}_n^V := \mathbf{h}_n \mathbf{h}_n^{\mathcal{T}}$$
 (1)

for all $n \in \mathcal{N}$, where we have introduced the measurement 141 vector $\mathbf{h}_n := \mathbf{e}_n$ with \mathbf{e}_n being the *n*th canonical vector in 142 \mathbb{R}^N . To express power injections as functions of \mathbf{v} , introduce 143 the bus admittance matrix $\mathbf{Y} = \mathbf{G} + j\mathbf{B} \in \mathbb{C}^N$, where \mathbf{G} and 144 $\mathbf{B} \in \mathbb{R}^{N \times N}$ are the real and imaginary parts of \mathbf{Y} [2], respectively. In rectangular coordinates, the active and reactive powers 146 P_n and Q_n injected into bus *n* can be expressed as 147

$$P_{n} = \Re(v_{n}) \sum_{n'=1}^{N} \left[\Re(v_{n'})G_{nn'} - \Im(v_{n})B_{nn'} \right] + \Im(v_{n}) \sum_{n'=1}^{N} \left[\Im(v_{n})G_{nn'} + \Re(v_{n})B_{nn'} \right]$$
(2)

$$Q_{n} = \Im(v_{n}) \sum_{n'=1}^{N} \left[\Re(v_{n})G_{nn'} - \Im(v_{n})B_{nn'} \right] - \Re(v_{n}) \sum_{n'=1}^{N} \left[\Im(v_{n})G_{nn'} + \Re(v_{n})B_{nn'} \right]$$
(3)

which can be compactly expressed as

 Δi

$$P_n = \mathbf{v}^{\mathcal{H}} \mathbf{H}_n^P \mathbf{v}, \text{ with } \mathbf{H}_n^P := \frac{\mathbf{Y}_n^{\mathcal{H}} + \mathbf{Y}_n}{2}$$
 (4a)

$$Q_n = \mathbf{v}^{\mathcal{H}} \mathbf{H}_n^Q \mathbf{v}, \text{ with } \mathbf{H}_n^Q := \frac{\mathbf{Y}_n^{\mathcal{H}} - \mathbf{Y}_n}{2j}$$
 (4b)

with $\mathbf{Y}_n := \mathbf{e}_n \mathbf{e}_n^T \mathbf{Y}$ for all $n \in \mathcal{N}$.

With regards to power flows, Kirchhoff's current law dictates 150 that the complex current over the line (n, n') at the "sending" 151 end is $i_{nn'} = y_{nn'}^s v_n + y_{nn'}(v_n - v'_n)$, where $y_{nn'}^s$ is the shunt 152 admittance at bus n' associated with the line (n, n'). The ac 153 power flow model, in conjunction with Ohm's law, further asserts that the complex power flowing over line (n, n') at the 155 "sending" end can be expressed as 156

$$S_{nn'} = P_{nn'} + jQ_{nn'} = v_n \overline{i}_{nn'}$$
$$= |v_n|^2 (\overline{y}_{nn'}^s + \overline{y}_{nn'}) - v_n \overline{v}_n \overline{y}_{nn'}.$$
(5)

It is worth pointing out that the complex power flow over line 157 $(n, n') \in \mathcal{L}$ at the "receiving" end is captured by that over line 158 $(n', n) \in \mathcal{L}$ at the "sending" end. Upon defining the following 159 matrices for all lines $(n, n') \in \mathcal{L}$: 160

$$\mathbf{Y}_{nn'} := (y_{nn'}^s + y_{nn'}) \mathbf{e}_n \mathbf{e}_n^{\mathcal{T}} - \overline{y}_{nn'} \mathbf{e}_{n'} \mathbf{e}_{n'}^{\mathcal{T}}$$
(6)

148

161 the active and reactive power flows at the "sending" terminal, 162 namely the real and imaginary parts of $S_{nn'}$ in (5) can be given 163 in a compact representation as

$$P_{nn'} = \mathbf{v}^{\mathcal{H}} \mathbf{H}_{nn'}^{P} \mathbf{v}, \text{ with } \mathbf{H}_{nn'}^{P} := \frac{\mathbf{Y}_{nn'}^{\mathcal{H}} + \mathbf{Y}_{nn'}}{2}$$
(7a)

$$Q_{nn'} = \mathbf{v}^{\mathcal{H}} \mathbf{H}_{nn'}^{Q} \mathbf{v}, \text{ with } \mathbf{H}_{nn'}^{Q} := \frac{\mathbf{Y}_{nn'}^{\mathcal{H}} - \mathbf{Y}_{nn'}}{2j}.$$
 (7b)

Having expressed all SCADA measurements as functions ofv, the PSSE problem can be presented next.

166 B. Power System State Estimation

167 In practice, the SCADA system measures a subset of the 168 system variables specified in (1), (4), and (7). Suppose now a 169 total of M such variables are measured, which are stacked up 170 to form the following $M \times 1$ measurement vector:

$$\mathbf{z} := \begin{bmatrix} \{\check{V}_n\}_{n \in \mathcal{N}_V}, \{\check{P}_n\}_{n \in \mathcal{N}_P}, \{\check{Q}_n\}_{n \in \mathcal{N}_Q} \\ \{\check{P}_{nn'}\}_{(n,n') \in \mathcal{L}_P}, \{\check{Q}_{nn'}\}_{(n,n') \in \mathcal{L}_Q} \end{bmatrix}^{\mathcal{T}} \in \mathbb{R}^M$$
(8)

where the check-marked terms represent possibly noisy observations of the corresponding error-free variables. The subsets $\mathcal{N}_V, \mathcal{N}_P, \mathcal{N}_Q \subseteq \mathcal{N}$, and $\mathcal{L}_P, \mathcal{L}_Q \subseteq \mathcal{L}$ specify the locations where meters are installed and the associated type of variables are measured. Succinctly, per *m*th measurement in z can be equivalently rewritten as

$$z_m := \mathbf{v}^{\mathcal{H}} \mathbf{H}_m \mathbf{v} + \epsilon_m \tag{9}$$

177 for all $m \in \{1, 2, ..., M\}$, where the terms $\epsilon_m \in \mathbb{R}$ capture 178 the metering errors and modeling inaccuracies, and the Hermi-179 tian measurement matrices $\mathbf{H}_m \in \mathbb{C}^{N \times N}$ can correspond to any 180 subset of the matrices defined in (1), (4), and (7). The critical 181 goal of PSSE is to obtain $\mathbf{v} \in \mathbb{C}^N$ based on the available data 182 $\{(z_m; \mathbf{H}_m)\}_{m=1}^M$.

183 Without loss of generality, adopting the LS error objective, 184 which coincides with the maximum likelihood criterion assum-185 ing additive white Gaussian noise, PSSE pursues problem¹

$$\underset{\mathbf{x}\in\mathbb{C}^{N}}{\operatorname{minimize}}\ \ell(\mathbf{x}) := \frac{1}{2M} \sum_{m=1}^{M} \left(z_{m} - \mathbf{x}^{\mathcal{H}} \mathbf{H}_{m} \mathbf{x} \right)^{2}.$$
(10)

Because of the quadratic terms inside the squares, the *quartic* function $\ell(\mathbf{x})$ is nonconvex, whose general instance is NP-hard [22]. Hence, it is computationally intractable to compute the LS or maximum likelihood estimate of **v** in general.

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190 C. Prior Contributions

Minimizing the nonlinear LS loss in (10), the Gauss–Newton method is the "workhorse" [2, Ch. 2]. Upon linearizing all quadratic terms $\mathbf{x}^{\mathcal{H}}\mathbf{H}_m\mathbf{x}$ around a given point using Taylor's expansion, the Gauss–Newton subsequently approximates the nonlinear LS fit in (10) using a linear one per iteration, and relies on its resultant minimizer to obtain the next iterate [3].

¹Throughout, \mathbf{v} is fixed for the actual system state, whereas \mathbf{x} is used for the optimization variable and the state estimate.

It typically converges in a few (≤ 10) iterations, very fast for 197 small- or medium-size problems. However, it is known that the 198 Gauss–Newton iterations for nonconvex LS are sensitive to the 199 initial point, and they may diverge in certain cases; see e.g., [5, 200 Ch. 5]. 201

On the other hand, several numerical polynomial-time SE al-202 gorithms have been pursued based on convex programming [6]. 203 By means of matrix lifting, such convex approaches start ex-204 pressing all quadratic measurements $\mathbf{x}^{\mathcal{H}} \mathbf{H}_m \mathbf{x}$ as linear functions 205 $Tr(\mathbf{H}_m \mathbf{X})$ of the rank-one matrix variable $\mathbf{X} := \mathbf{x}\mathbf{x}^{\mathcal{H}} \succeq \mathbf{0}$. 206 Upon discarding the nonconvex rank constraint, the nonlinear 207 LS in (10) boils down to (or can be converted into) a convex 208 SDP. In terms of computational efficiency, such convex schemes 209 entail solving for an $N \times N$ positive semidefinite matrix from 210 M SDP constraints, whose worst case computational complex-211 ity is $\mathcal{O}(M^4 N^{1/2} \log(1/\epsilon))$ for any given solution accuracy 212 $\epsilon > 0$ [23, Sec. 6.6.2]. This complexity and the resultant storage 213 requirement evidently do not scale nicely to the increasingly 214 interconnected large power networks. 215

III. RANK-ONE MEASUREMENT APPROACH 216

Drawing from advances in nonconvex optimization, this sec-217 tion presents a new framework for scalable, accurate, and ro-218 bust PSSE. Specifically, our proposed approach reformulates 219 the rank-two power flows in (7) into rank-one quadratic mea-220 surements (namely with corresponding measurement matrices 221 having rank one), followed by two efficient algorithms for min-222 imizing the ℓ_1 -based misfit of the transformed measurements. 223 In contrast to previous SE approaches that minimize the quar-224 *tic* polynomial in (10), a novel *quadratic* objective functional 225 is obtained and subsequently minimized. With more complicated algebraic manipulations, the power injections can also 227 be accounted for in our proposed framework. In the presence 228 of additive noise, near-optimal statistical performance of the 229 developed approach is numerically demonstrated. 230

A. Measurement Transformation

In this paper, we focus on the following types of measure-232 ments: first, $\{|v_m|\}_{m=1}^N$ the voltage magnitudes at all buses, 233 and second, $\{P_{nn'}\}_{(m,n)\in\mathcal{L}_P}$ and/or $\{Q_{nn'}\}_{(m,n)\in\mathcal{L}_Q}$ the ac-234 tive and/or reactive power flows on a selected subset of lines, 235 namely \mathcal{L}_P , $\mathcal{L}_Q \subseteq \mathcal{L}$. Consider first the noise-free case, where 236 all available measurements can be described as 237

$$z_m = \mathbf{v}^{\mathcal{H}} \mathbf{H}_m \mathbf{v} \ \forall m = 1, \dots, M.$$
(11)

Without loss of generality, let the first N measurements be the 238 squared voltage magnitudes at the N buses, namely $z_m = |v_m|^2$ 239 for m = 1, 2, ..., N, and the remaining ones be the (active or 240 reactive) power flows. It is clear from (1) that the squared voltage 241 magnitude measurements are given by 242

$$z_m = \left| \mathbf{h}_m^{\mathcal{H}} \mathbf{v} \right|^2 \quad \forall m = 1, \dots, N \tag{12}$$

whose corresponding measurement matrices have rank one; that 243 is, $\{\mathbf{H}_m = \mathbf{h}_m \mathbf{h}_m^{\mathcal{H}}\}_{m=1}^N$. 244

Now let us consider the power flow data pairs $(P_{nn'}; \mathbf{H}_{nn'}^P)$ 245 and $(Q_{nn'}; \mathbf{H}_{nn'}^Q)$ in (7). Upon substituting $\mathbf{Y}_{nn'}$ in (6) into (7), 246

247 one can rewrite for all lines $(n, n') \in \mathcal{L}_P$

$$\mathbf{H}_{nn'}^{P} = \frac{1}{2} \left(\alpha_{nn'}^{P} \mathbf{e}_{n} \mathbf{e}_{n}^{T} - \beta_{nn'}^{P} \mathbf{e}_{n} \mathbf{e}_{n'}^{T} - \bar{\beta}_{nn'}^{P} \mathbf{e}_{n'} \mathbf{e}_{n'}^{T} \mathbf{e}_{n'}^{T}$$

248 and similarly for all lines $(n, n') \in \mathcal{L}_Q$

$$\mathbf{H}_{nn'}^{Q} = \frac{1}{2} \left(\alpha_{nn'}^{Q} \mathbf{e}_{n} \mathbf{e}_{n}^{T} - \beta_{nn'}^{Q} \mathbf{e}_{n} \mathbf{e}_{n'}^{T} - \bar{\beta}_{nn'}^{Q} \mathbf{e}_{n'} \mathbf{e}_{n'}^{T} \mathbf{e}_{n'}^{T}$$

249 where the four coefficients are given by

$$\alpha_{nn'}^P := 2\Re(y_{nn'}^s + y_{nn'}), \qquad \beta_{nn'}^P := y_{nn'}$$
(15a)

$$\alpha_{nn'}^Q := -2\Im(y_{nn'}^s + y_{nn'}), \qquad \beta_{nn'}^Q := jy_{nn'}.$$
(15b)

It is worth pointing out that $\alpha_{nn'}^P$ and $\alpha_{nn'}^Q$ are both real, whereas $\beta_{nn'}^P$ and $\beta_{nn'}^Q$ are in general complex.

It is self-evident from (13) and (14) that each of the active or reactive power flow measurement matrices $\{\mathbf{H}_{nn'}^{P}\}_{(n,n')\in\mathcal{L}_{P}}$ or $\{\mathbf{H}_{nn'}^{Q}\}_{(n,n')\in\mathcal{L}_{Q}}$ is of *at most* rank-two, with three nonzero matrix entries at the (n, n)-, (n, n')-, and (n', n)th positions depending on whether the involved coefficients are nonzero or not. Yet, these power flow matrices are generally indefinite. Upon neglecting the superscripts for notational brevity, both $\mathbf{H}_{nn'}^{P}$ and $\mathbf{H}_{nn'}^{Q}$ are in the following form:

$$\mathbf{H}_{nn'} = \frac{1}{2} \left(\alpha_{nn'} \mathbf{e}_n \mathbf{e}_n^{\mathcal{T}} - \beta_{nn'} \mathbf{e}_n \mathbf{e}_{n'}^{\mathcal{T}} - \bar{\beta}_{nn'} \mathbf{e}_{n'} \mathbf{e}_n^{\mathcal{T}} \right).$$
(16)

We establish the following result for power flow data $\{z_m\}$. *Proposition 1 (Rank-one measurements):* Suppose that the voltage magnitudes are measured at all buses. Then, one can construct equivalently a new measurement \check{z}_m for each power flow z_m so that its corresponding measurement matrix $\check{\mathbf{H}}_{nn'}$ is of rank one, and positive semidefinite.

Proof: Depending on whether the coefficient $\alpha_{nn'}$ or $\beta_{nn'}$ is zero or not, we discuss separately the ensuing three cases:

268 c1) $\alpha_{nn'} \neq 0, \ \beta_{nn'} \neq 0;$ 269 c2) $\alpha_{nn'} = 0, \ \beta_{nn'} \neq 0;$

270 c3) $\alpha_{nn'} = 0, \beta_{nn'} \neq 0$

Q3

271 where the trivial case of $\alpha_{nn'} = \beta_{nn'} = 0$ is excluded because 272 all SCADA measurements are assumed nonzero. The goal of 273 the following section is to transform all power flow measure-274 ment matrices into *rank-one (symmetric) positive semidefinite* 275 matrices. The three cases are individually discussed next.

276 Consider first the case c1). When both $\alpha_{nn'}$ and $\beta_{nn'}$ are 277 nonzero, the matrix $\mathbf{H}_{nn'}$ has exactly rank two with three 278 nonzero entries. Its two nonzero eigenvalues can be obtained 279 by solving the quadratic equation

$$\lambda \left(\lambda - \frac{\alpha_{nn'}}{2} \right) - \frac{|\beta_{nn'}|^2}{4} = 0$$

which is derived by setting the determinant of $(\lambda \mathbf{I}_N - \mathbf{H}_{nn'})$ to zero. Its closed-form solutions are given by

$$\lambda_1 = \frac{\alpha_{nn'} + \sqrt{\alpha_{nn'}^2 + 4|\beta_{nn'}|^2}}{4} > 0$$

$$\lambda_2 = \frac{\alpha_{nn'} - \sqrt{\alpha_{nn'}^2 + 4|\beta_{nn'}|^2}}{4} < 0.$$

Let $\mathbf{u}_1, \mathbf{u}_2 \in \mathbb{C}^N$ be the unit eigenvectors of $\mathbf{H}_{nn'}$ associated with the eigenvalues λ_1, λ_2 , respectively. Hence, one can write $\mathbf{H}_{nn'} := \lambda_1 \mathbf{u}_1 \mathbf{u}_1^{\mathcal{H}} + \lambda_2 \mathbf{u}_2 \mathbf{u}_2^{\mathcal{H}}$. To obtain a rank-one positive semidefinite matrix, the first attempt would be to compensate for the negative eigenvalue λ_2 and make it zero. This is tantamount to adding $-\lambda_2 \mathbf{u}_2 \mathbf{u}_2^{\mathcal{H}}$ to $\mathbf{H}_{nn'}$, and accordingly adding $-\lambda_2 \mathbf{v}^{\mathcal{H}} (\mathbf{u}_2 \mathbf{u}_2^{\mathcal{H}}) \mathbf{v}$ to the measurement $z_{nn'}$; that is 288

$$\mathbf{\check{H}}_{nn'} := \mathbf{H}_{nn'} - \lambda_2 \mathbf{u}_2 \mathbf{u}_2^{\mathcal{H}}$$
(17a)

$$\check{z}_{nn'} := z_{nn'} - \lambda_2 \mathbf{v}^{\mathcal{H}} (\mathbf{u}_2 \mathbf{u}_2^{\mathcal{H}}) \mathbf{v}$$
(17b)

in which the transformed measurement matrix $\hat{\mathbf{H}}_{nn'}$ is rank-one 289 and symmetric positive semidefinite, and $\check{z}_{nn'}$ is the resultant 290 transformed measurement. To realize this however, entails, eval-291 uating the term $\mathbf{v}^{\mathcal{H}}(\mathbf{u}_2 \mathbf{u}_2^{\mathcal{H}})\mathbf{v}$, which requires knowledge of the 292 true state vector \mathbf{v} . This procedure is thus not feasible, and one 293 has to develop a new twist to bypass this hurdle. 294

Recall from our working assumption that we have access to 295 all squared voltage magnitudes $\{|v_n|^2\}_{n=1}^N$. Based on this fact, 296 we show in the following that it is sufficient to add a matrix of 297 the form $(\delta_{nn'}/2)\mathbf{e}_{n'}\mathbf{e}_{n'}^T$ to $\mathbf{H}_{nn'}$ such that the resulting sum, 298 denoted by 299

$$\mathbf{\check{H}}_{nn'} := \mathbf{H}_{nn'} + (\delta_{nn'}/2)\mathbf{e}_{n'}\mathbf{e}_{n'}^T$$
(18)

can be rendered rank-one. Here, $\delta_{nn'}$ is an unknown coefficient 300 to be determined next. Toward this end, setting the determinant 301 of $(\lambda \mathbf{I}_N - \check{\mathbf{H}}_{nn'})$ to zero leads to 302

$$\left(\lambda - \frac{\alpha_{nn'}}{2}\right) \left(\lambda - \frac{\delta_{nn'}}{2}\right) - \frac{|\beta_{nn'}|^2}{4} = 0.$$
(19)

To yield a rank-one matrix $\hat{\mathbf{H}}_{nn'}$, it is sufficient for the 303 quadratic equation (19) to have exactly one nonzero solution. By 304 basic linear algebra, this is equivalent to having a zero constant 305 term in (19), giving rise to 306

$$\alpha_{nn'}\delta_{nn'} - |\beta_{nn'}|^2 = 0$$

or alternatively

 $\delta_{nn'} := |\beta_{nn'}|^2 / \alpha_{nn'}.$

It can be verified that the transformed measurement matrix 308

$$\check{\mathbf{H}}_{nn'} := \mathbf{H}_{nn'} + (|\beta_{nn'}|^2 / (2\alpha_{nn'})) \mathbf{e}_{n'} \mathbf{e}_{n'}^{\mathcal{T}}$$
(20)

is rank-one. In addition, if $\alpha_{nn'} > 0$, then $\dot{\mathbf{H}}_{nn'}$ is positive 309 semidefinite. Therefore, one can write 310

$$\dot{\mathbf{H}}_{nn'} := \mathbf{h}_{nn'} \mathbf{h}_{nn'}^{\mathcal{H}} \tag{21}$$

with the equivalent measurement vector being

$$\mathbf{h}_{nn'} := \sqrt{\frac{\alpha_{nn'}}{2}} \mathbf{e}_n + \frac{\bar{\beta}_{nn'}}{\sqrt{2\alpha_{nn'}}} \mathbf{e}_{n'}.$$
 (22)

The transformed measurement $\check{z}_{nn'}$ corresponding to $\hat{\mathbf{H}}_{nn'}$ can 312 be given by 313

$$\check{z}_{nn'} := \mathbf{v}^{\mathcal{H}} \check{\mathbf{H}}_{nn'} \mathbf{v} = \mathbf{v}^{\mathcal{H}} \mathbf{H}_{nn'} \mathbf{v} + \mathbf{v}^{\mathcal{H}} \left(\delta_{nn'} \mathbf{e}_{n'} \mathbf{e}_{n'}^{\mathcal{T}} / 2 \right) \mathbf{v}
= z_{nn'} + \left(|\beta_{nn'}|^2 / 2\alpha_{nn'} \right) |v_{n'}|^2 \quad (23)$$

307

for which the required quantity $(|\beta_{nn'}|^2/2\alpha_{nn'})|v_{n'}|^2$ is available, or can be obtained as long as $|v_n|^2$ is available.

316 If $\alpha_{nn'} < 0$, one can instead define

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$$\check{\mathbf{H}}_{nn'} := -\mathbf{H}_{nn'} + (|\beta_{nn'}|^2/(2\alpha_{nn'}))\mathbf{e}_{n'}\mathbf{e}_{n'}^{\mathcal{T}}$$
(24)

317 and write the equivalent measurement vector as

$$\mathbf{h}_{nn'} := \sqrt{\frac{-\alpha_{nn'}}{2}} \mathbf{e}_n + \frac{\beta_{nn'}}{\sqrt{-2\alpha_{nn'}}} \mathbf{e}_{n'}$$
(25)

318 for which the corresponding measurement becomes

$$\begin{aligned} \tilde{\mathbf{z}}_{nn'} &:= \mathbf{v}^{\mathcal{H}} \mathbf{\dot{H}}_{nn'} \mathbf{v} = -\mathbf{v}^{\mathcal{H}} \mathbf{H}_{nn'} \mathbf{v} - \mathbf{v}^{\mathcal{H}} \left(\delta_{nn'} \mathbf{e}_{n'} \mathbf{e}_{n'}^2 / 2 \right) \mathbf{v} \\ &= -z_{nn'} - \left(|\beta_{nn'}|^2 / 2\alpha_{nn'} \right) |v_n|^2. \end{aligned}$$
(26)

19 Likewise, the quantity $-(|\beta_{nn'}|^2/2\alpha_{nn'})|v_n|^2$ is also available under our working assumption.

Let us now focus on the case c2). To transform $\mathbf{H}_{nn'}$ into a rank-one positive semidefinite matrix, it is sufficient to add a matrix of the form $(\gamma_{nn'}/2)\mathbf{e}_n\mathbf{e}_n^T + (\delta_{nn'}/2)\mathbf{e}_{n'}\mathbf{e}_{n'}^T$, to yield

$$\check{\mathbf{H}}_{nn'} := \mathbf{H}_{nn'} + (\gamma_{nn'}/2)\mathbf{e}_n \mathbf{e}_n^{\mathcal{T}} + (\delta_{nn'}/2)\mathbf{e}_{n'} \mathbf{e}_{n'}^{\mathcal{T}}$$
(27)

for some coefficients $\gamma_{nn'} > 0$ and $\delta_{nn'} > 0$ to be determined. Similar to the discussion for case c1), to find $\gamma_{nn'}$ and $\delta_{nn'}$, one sets the determinant of $(\lambda \mathbf{I}_N - \check{\mathbf{H}}_{nn'})$ to zero, leading to

$$\left(\lambda - \frac{\gamma_{nn'}}{2}\right)\left(\lambda - \frac{\delta_{nn'}}{2}\right) - \frac{|\beta_{nn'}|^2}{4} =$$

327 The fact that $\mathbf{\hat{H}}_{nn'}$ is rank-one implies that

$$\gamma_{nn'}\delta_{nn'} - |\beta_{nn'}|^2 = 0.$$

328 Without loss of generality, one can take

$$\gamma_{nn'} := 1$$
, and $\delta_{nn'} := |\beta_{nn'}|^2$

329 and $\mathbf{H}_{nn'}$ in (27) becomes rank-one and can be written as

$$\check{\mathbf{H}}_{nn'} := \mathbf{h}_{nn'} \mathbf{h}_{nn'}^{\mathcal{H}}$$
(28)

0.

330 with

$$\mathbf{h}_{nn'} := \frac{1}{\sqrt{2}} \mathbf{e}_n - \frac{\beta_{nn'}}{\sqrt{2}} \mathbf{e}_{n'}.$$
 (29)

331 The transformed measurement associated with $\hat{\mathbf{H}}_{nn'}$ can be 332 found as follows:

$$\check{z}_{nn'} := \mathbf{v}^{\mathcal{H}} \mathbf{H}_{nn'} \mathbf{v}$$

$$= \mathbf{v}^{\mathcal{H}} \mathbf{H}_{nn'} \mathbf{v} + \mathbf{v}^{\mathcal{H}} \left(\beta_{nn'} \mathbf{e}_n \mathbf{e}_n^{\mathcal{T}} / 2 + \delta_{nn'} \mathbf{e}_{n'} \mathbf{e}_{n'} / 2 \right) \mathbf{v}$$

$$= z_{nn'} + (1/2) |v_n|^2 + (|\beta_{nn'}|^2 / 2) |v_{n'}|^2$$
(30)

for which the required quantities $|v_n|^2$ and $|\beta_{nn'}|^2 |v_{n'}|^2$ can be computed when $|v_n|^2$ and $|v_{n'}|^2$ are available.

Let us now turn to the last case c3). Since $\beta_{nn'} = \overline{\beta}_{nn'} = 0$, one can write $\mathbf{H}_{nn'} = \alpha_{nn'} \mathbf{e}_n \mathbf{e}_n^T / 2$, which is already rank-one. To make it positive semidefinite, it suffices to take the absolute value, and define

$$\check{\mathbf{H}}_{nn'} := \mathbf{h}_{nn'} \mathbf{h}_{nn'}^{\mathcal{H}} = \left(\sqrt{\frac{|\alpha_{nn'}|}{2}} \mathbf{e}_n\right) \left(\sqrt{\frac{|\alpha_{nn'}|}{2}} \mathbf{e}_n\right)^T.$$
(31)

The transformed measurement is given by

$$\check{z}_{nn'} = |z_{nn'}|. \tag{32}$$

To summarize, for any active or reactive power flow data 340 $(z_{nn'}; \mathbf{H}_{nn'})$, we have developed a strategy to obtain a new 341 measurement pair $(\check{z}_{nn'}; \check{\mathbf{H}}_{nn'})$, in which $\check{\mathbf{H}}_{nn'}$ becomes rank- 342 one and positive semidefinite. Specifically, this is accomplished 343 through steps in (21)–(32), by depending upon the values of 344 coefficients $\alpha_{nn'}$ and $\beta_{nn'}$, provided that the voltage magnitudes 345 at all buses are available.

The assumption on full voltage measurements can be relaxed. 347 Indeed, it is possible to build up rank-one measurements via 348 linear combinations, so long as two of the following SCADA 349 quantities $\{|v_n|, |v_{n'}|, p_{nn'}, p_{n'n}, q_{nn'}, q_{n'n}\}$ are measured on 350 every line $(n, n') \in \mathcal{L}$. In a nutshell, one can readily rewrite all 351 measurements as intensities of some known and deterministic 352 linear transforms of the state vector, namely 353

$$\check{z}_m = \left| \mathbf{h}_m^{\mathcal{H}} \mathbf{v} \right|^2 \quad \forall m = 1, \dots, M$$
(33)

where the measurement vectors $\mathbf{h}_m \in \mathbb{C}^N$ are given in (1), (22), 354 (25), (29), and (31), whereas the corresponding transformed 355 measurements $\check{z}_m > 0$ are defined in (1), (23), (26), (30), and 356 (32). Moreover, all vectors \mathbf{h}_m are highly sparse, each having 357 at most two nonzero entries independent of the system size N. 358 This feature can be carefully exploited to endow the iterative 359 PSSE solvers with computational efficiency and scalability. 360

IV. PROX-LINEAR SE SOLVERS

In general, given a set of (consistent) quadratic equations, 362 there may exist multiple solutions even after excluding triv-363 ial ambiguities. In the context of phase retrieval, in which the 364 measurement matrices are rank-one positive semidefinite (i.e., 365 $\mathbf{H}_m = \mathbf{h}_m \mathbf{h}_m^{\mathcal{H}}$), a number $M \ge 4N - 4$ of random quadratic 366 equations suffice for uniqueness of the solution [24]. In the 367 power systems literature though, it remains an open question 368 that how many quadratic measurements as in (1), (4), and (7)369 are required for the uniqueness of PSSE solution. For concrete-370 ness, this contribution assumes that a large enough number of 371 measurements are available, and they collectively determine a 372 unique solution, namely the underlying true system state. 373

Leveraging the rank-one measurement model, this section 374 presents two algorithms for scalable and exact power system 375 state recovery based on nonconvex optimization. Specifically, 376 we focus on the ℓ_1 -loss to fit the intensity measurements \check{z}_m 377 in (33) instead of z_m in (11). Despite the nonconvexity and 378 nonsmoothness of the resulting loss function, we first develop 379 a deterministic prox-linear algorithm. When the initialization is 380 sufficiently close to the underlying true voltage state vector, and 381 the loss function satisfies a certain local "stability condition," 382 we show that our first deterministic approach recovers the true 383 voltage vector at a quadratic convergence rate. It entails solv-384 ing a quadratic program per iteration, for which off-the-shelf 385 convex programming toolboxes are widely available, but the re-386 sulting complexity does not scale well with the system size. To 387 endow the algorithm with scalability, a stochastic generalization 388 is pursued, which processes a single measurement per iteration. 389

339

It is well known in statistics and power systems literature that ℓ_1 -based loss functions yield median-based estimators, and they can cope with gross errors in the measurements \check{z}_m in a relatively benign way [21]. This prompts us to consider the ℓ_1 -loss [i.e., least-absolute-value (LAV)] formulation

$$\underset{\mathbf{x}\in\mathbb{C}^{N}}{\operatorname{minimize}}\ \ell(\mathbf{x}) := \frac{1}{M} \sum_{m=1}^{M} \left| \check{z}_{m} - |\mathbf{h}_{m}^{\mathcal{H}} \mathbf{x}|^{2} \right|.$$
(34)

Evidently, this loss function is nonsmooth and nonconvex, which is not even locally convex near v. This can be understood from the scalar case $\ell(x) = |1 - x^2|$. As such, it is unclear how to efficiently minimize such functions.

However, the function exhibits several appealing structural 399 properties that we explore in the following to develop iterative 400 algorithms to locally solve the problem (34) efficiently [21], 401 [25], [26]. To this end, consider first expressing the loss func-402 tion as the composition $\ell(\mathbf{x}) := c(\mathbf{s}(\mathbf{x}))$, of the convex func-403 tion $c(\cdot) := \|\cdot\|_1$ and the smooth one $\mathbf{s}(\mathbf{x}) := \frac{1}{M} (\check{\mathbf{z}} - |\mathbf{H}\mathbf{x}|^2),$ 404 where $\check{\mathbf{z}} := [\check{z}_1 \cdots \check{z}_M]^T$, $\mathbf{H} := [\mathbf{h}_1 \cdots \mathbf{h}_M]^{\mathcal{H}}$, and the mod-405 ulus operator $|\cdot|$ is understood elementwise when applied to 406 a vector. Such a compositional structure lends itself nicely to 407 iterative procedures that are referred to as proximal-linear (prox-408 linear) algorithms, which we are described in detail ahead. 409

410 Consider a real-valued $\mathbf{x} \in \mathbb{R}^N$ for now. The (deterministic) 411 prox-linear method for minimizing $\ell(\mathbf{x}) = c(\mathbf{s}(\mathbf{x}))$ is to lin-412 earize s only, followed by successively minimizing a sequence 413 of locally regularized models. Specifically, starting with some 414 initialization $\mathbf{x}_0 \in \mathbb{R}^N$, the prox-linear method defines a local 415 "linearization" of ℓ around the current iterate $\mathbf{x}_t \in \mathbb{R}^N$ as

$$\ell_{\mathbf{x}_t}(\mathbf{x}) := c \big(\mathbf{s}(\mathbf{x}_t) + \nabla^T \mathbf{s}(\mathbf{x}_t)(\mathbf{x} - \mathbf{x}_t) \big)$$
(35)

416 with $\nabla \mathbf{s}(\mathbf{x}_t) \in \mathbb{R}^{N \times M}$ representing the Jacobian matrix of s 417 evaluated at point \mathbf{x}_t ; and subsequently, it proceeds inductively 418 to obtain iterates $\mathbf{x}_1, \mathbf{x}_2, ...$ by minimizing the quadratically reg-419 ularized models [25], [26]

$$\mathbf{x}_{t+1} = \arg\min_{\mathbf{x}\in\mathbb{R}^N} \left\{ \ell_{\mathbf{x}_t}(\mathbf{x}) + \frac{1}{2\mu_t} \|\mathbf{x} - \mathbf{x}_t\|_2^2 \right\}$$
(36)

where $\mu_t > 0$ is a step size that can be fixed *a priori* to some 420 constant, or be determined "on-the-fly" through a line search 421 [25], [26]. Furthermore, observing that the linearization $\ell_{\mathbf{x}_{t}}(\mathbf{x})$ 422 is convex in x, so problem (36) is convex in x as well. It has 423 been shown in [26] that when c is L-Lipschitz and ∇s is κ -424 Lipschitz, choosing any step size $0 < \mu < \frac{1}{\kappa L}$ guarantees that 425 the algorithm (36) is a descent method; that is, the iterates $\{\mathbf{x}_k\}$ 426 monotonically decrease the function value of $\ell(\mathbf{x})$; and finds an 427 (approximate) stationary point of (34). 428

Nevertheless, the PSSE problem (34) involves optimization 429 over complex-valued variables in $\mathbf{x} \in \mathbb{C}^N$. It can be checked 430 that the functions ℓ and s do not satisfy the Cauchy–Riemann 431 (CR) equations; see e.g., [27, Th. 7.2] for the definition of CR 432 equations. Hence, functions ℓ and s are not holomorphic (i.e., 433 complex-differentiable) in x. As such, the "linearization," or the 434 first-order Taylor's expansion of $\mathbf{s}(\mathbf{x})$ in $\mathbf{x} \in \mathbb{C}^N$ alone [cf. (35)] 435 does not exist. To address this challenge, we invoke Wirtinger's 436 calculus to generalize prox-linear algorithms to optimization 437

Algorithm 1: Deterministic Prox-linear SE Solver.

- 1: Input data $\{(z_m, \mathbf{H}_m)\}_{m=1}^M$, step size $\mu > 0$, initialization $\mathbf{v}_0 \in \mathbb{C}^N$, solution accuracy $\epsilon > 0$, and set t = 0.
- 2: Prepare the power flow data {(z_m, H_m)}^M_{m=N+1} according to (1), (22)–(23), (25)–(26), (29)–(30), and (31)–(32) to obtain {(ž_m, h_m)}^M_{m=N+1} based on {z_m = |v_m|²}^N_{m=1}.
 3: Repeat
 4: Evaluate a_{m,t} and b_{m,t} in (38).
- 5: Solve (37) to yield \mathbf{x}_{t+1} .
- 6: t = t + 1.
- 7: Until $\|\mathbf{x}_t \mathbf{x}_{t-1}\|_2 \leq \epsilon \sqrt{N}$.
- 8: **Return** \mathbf{x}_t .

over complex-valued arguments in the sequel. Please refer to 438 [28] for basics of Wirtinger's calculus. 439

A. Deterministic Prox-Linear SE Solver

Our first deterministic prox-linear approach to (34) is simply 441 stated: begin with initialization $\mathbf{x}_0 := \mathbf{1} \in \mathbb{R}^N$, and proceed 442 by successively minimizing quadratically regularized functions 443 around the current iterate $\mathbf{x}_t \in \mathbb{C}^N$ to yield the next iterate 444

$$\mathbf{x}_{t+1} = \arg\min_{\mathbf{x}\in\mathbb{C}^N} \frac{1}{M} \sum_{m=1}^M \left| b_{m,t} - 2\Re (\mathbf{a}_{m,t}^{\mathcal{H}} \mathbf{x}) \right| + \frac{1}{2\mu_t} \|\mathbf{x} - \mathbf{x}_t\|_2^2$$
(37)

where the term $b_{m,t} - 2\Re(\mathbf{a}_{m,t}^{\mathcal{H}}\mathbf{x})$ can be interpreted as the firstorder Taylor's approximation of the nonholomophic function 446 $\check{z}_m - |\mathbf{h}_m^{\mathcal{H}}\mathbf{x}|^2$ at \mathbf{x}_t based upon the Wirtinger derivatives; see 447 Appendix A for the rigorous derivation. The coefficients $\mathbf{a}_{m,t}$ 448 and $b_{m,t}$ are given by 449

$$\mathbf{a}_{m,t} := \left(\mathbf{h}_m^{\mathcal{H}} \mathbf{x}_t\right) \mathbf{h}_m \tag{38a}$$

$$b_{m,t} := \check{z}_m + \left| \mathbf{h}_m^{\mathcal{H}} \mathbf{x}_t \right|^2.$$
(38b)

Observe that the problem (37) to be tackled per iteration of 450 our deterministic prox-linear SE solver is a convex quadratic 451 program, which can be efficiently solved with standard convex programming methods. Under appropriate conditions, our 453 scheme converges quadratically fast to the true state vector 454 \mathbf{v} , meaning that we have to solve only about $\log_2 \log_2(1/\epsilon)$ 455 such quadratic programs to obtain an estimate \mathbf{x} of \mathbf{v} satisfying 456 dist $(\mathbf{x}, \mathbf{v}) \leq \epsilon ||\mathbf{v}||_2$. As will be corroborated by our numerical 457 tests in Section VI, this boils down to $5 \sim 8$ convex quadratic 458 programs in practice. Moreover, our approach applies both in 459 the noiseless setting, and when a constant (random) portion of 460 the measurements are even adversarially corrupted.

For implementation, our deterministic prox-linear solver is 462 summarized in Algorithm 1. Regarding computational complexity, preparing the data in Step 2 can be performed within 464 $\mathcal{O}(M)$ operations. Exploiting the sparsity of \mathbf{h}_m 's, evaluating 465 the coefficients $\{(b_{m,t}, \mathbf{a}_{m,t})\}_{m=1}^M$ in Step 5 can also be done 466 with $\mathcal{O}(M)$ operations. The overall complexity of Algorithm 1 467

is indeed dominated by solving the quadratic program of (37) in 468 Step 6. With standard convex programming solvers, the resultant 469 complexity is often $\mathcal{O}(MN^2)$. Iterative procedures depending 470 471 on the alternating direction method of multipliers can reduce this number to $\mathcal{O}(MN\log(1/\epsilon))$ [21], [29]. The latter complexity, 472 however, may still become unfavorable for large-size power net-473 works. Furthermore, even though $\{\mathbf{h}_m\}_{m=1}^M$ have at most two 474 nonzero entries, this property cannot be fully exploited to speed 475 up computations for solving the quadratic program of (37). To 476 477 address these issues, we advocate a stochastic alternative of (37)ahead for solving problem (34). 478

479 B. Stochastic Prox-Linear SE Solver

The stochastic prox-linear method deals with a single measurement per iteration. Initialized with \mathbf{x}_0 , our (stochastic) proxlinear SE solver operates by first sampling uniformly a loss function via randomly picking $m_t \in \{1, 2, ..., M\}$, and relies on minimizing its quadratically regularized "linearization" around \mathbf{x}_t to yield \mathbf{x}_{t+1} [30]; that is, define inductively for t = 0, 1, 2, ... that

$$\mathbf{x}_{t+1} = \arg\min_{\mathbf{x}\in\mathbb{C}^N} \left| b_{m_t,t} - 2\Re \left(\mathbf{a}_{m_t,t}^{\mathcal{H}} \mathbf{x} \right) \right| + \frac{1}{2\mu_t} \|\mathbf{x} - \mathbf{x}_t\|_2^2$$
(39)

where the coefficients $b_{m_t,t}$ and $\mathbf{a}_{m_t,t}$ are given in (38), with $b_{m_t,t} - 2\Re(\mathbf{a}_{m_t,t}^{\mathcal{H}}\mathbf{x})$ being the first-order Taylor's expansion of the m_t th error function $\check{z}_{m_t} - |\mathbf{h}_{m_t}^{\mathcal{H}}\mathbf{x}|^2$ around \mathbf{x}_t . Evidently, problem (39) is again a quadratic program too. Compared with the first quadratic program in (36), fortunately, the solution to (39) can be found in simple closed form.

To that end, we invoke an earlier result in [29, Prop. 3], which is included in Appendix B for completeness. Upon defining $\mathbf{w} := \mathbf{x} - \mathbf{x}_t$, $\mathbf{a} := \mathbf{a}_{m_t,t}$, and $b := b_{m_t,t} - 2\Re(\mathbf{a}_{m_t,t}^{\mathcal{H}}\mathbf{x}_t)$ in Proposition 2, one can readily find the solution to (39) as follows:

$$\mathbf{x}_{t+1} = \mathbf{x}_t + \operatorname{proj}_{\mu_t} \left(\frac{b_{m_t,t} - 2\Re(\mathbf{a}_{m_t,t}^{\mathcal{H}} \mathbf{x}_t)}{\|\mathbf{a}_{m_t,t}\|_2^2} \right) \mathbf{a}_{m_t,t}.$$

498 Substituting \mathbf{a}_{m_t} and b_{m_t} of (38) into the last equality leads to 499 our stochastic prox-linear SE solver

$$\mathbf{x}_{t+1} = \mathbf{x}_t + \operatorname{proj}_{\mu_t} \left(\frac{\check{z}_{m_t} - |\mathbf{h}_{m_t}^{\mathcal{H}} \mathbf{x}_t|^2}{4 |\mathbf{h}_{m_t}^{\mathcal{H}} \mathbf{x}_t|^2 \cdot ||\mathbf{h}_{m_t}||_2^2} \right) \cdot 2 (\mathbf{h}_{m_t}^{\mathcal{H}} \mathbf{x}_t) \mathbf{h}_{m_t}.$$
(40)

We summarize our second (stochastic) prox-linear SE solver 500 in Algorithm 2 for further reference. In terms of computational 501 complexity, we report the exact number of complex scalar op-502 erations (e.g., additions, multiplications) needed per stochas-503 tic prox-linear SE iteration of (40) next. Relying on whether 504 \mathbf{h}_{m_t} has 1 or 2 nonzero entries, the following statements hold 505 true. If \mathbf{h}_{m_t} has 1 (2) nonzero entries, evaluating $|\mathbf{h}_{m_t}^{\mathcal{H}} \mathbf{x}_t|^2$ re-506 quires 2 (4) operations, and $\|\mathbf{h}_{m_t}\|_2^2$ requires 1 (3) operations, 507 plus another 5 operations for the remaining, all summing to 508 a total of 8 (12) operations. In other words, per iteration of 509 Algorithm 2 (cf. Steps 4-6) must perform only 12 complex 510 scalar operations or so. Interestingly, this per-iteration com-511 plexity of $\mathcal{O}(1)$ holds regardless of the power network under 512

Algorithm 2: Stochastic Prox-linear SE Solver.

- 1: Input data $\{(z_m, \mathbf{H}_m)\}_{m=1}^M$, step size $\mu > 0$, initialization $\mathbf{v}_0 \in \mathbb{C}^N$, solution accuracy $\epsilon > 0$, and set t = 0.
- 2: Prepare the power flow data {(z_m, H_m)}^M_{m=N+1} according to (1), (22)–(23), (25)–(26), (29)–(30), and (31)–(32) to obtain {(ž_m, h_m)}^M_{m=N+1} based on {z_m = |v_m|²}^N_{m=1}.
 3: Repeat
 4: Draw m_t ∈ {1, 2, ..., M} uniformly at random.
- 5: **Evaluate** Evaluate $\mathbf{a}_{m_t,t}$ and $b_{m_t,t}$ in (38).
- 6: **Update** x_{t+1} via (40).
- 7: t = t + 1.

8: Until
$$\|\mathbf{x}_t - \mathbf{x}_{t-1}\|_2 \leq \epsilon \sqrt{N}$$

9: **Return** \mathbf{x}_t .

investigation, or more precisely, the system size N. It is selfevident that this $\mathcal{O}(1)$ per-iteration complexity scales nicely to large- and even massive-size power networks. 515

V. CONVERGENCE ANALYSIS AND EXACT RECOVERY 516

In this section, we begin our development by providing convergence guarantees for the proposed prox-linear SE solvers. 518 For concreteness, we will focus on the deterministic prox-linear 519 Algorithm 1, whereas convergence can be also established for 520 Algorithm 2 in a probabilistic sense. Interested readers are referred to [30]. Under certain conditions on the loss function ℓ , 522 exact recovery results are also established. 523

Recall our loss function $c(\mathbf{s}(\mathbf{x})) = \frac{1}{M} \|\mathbf{\check{z}} - |\mathbf{H}\mathbf{x}|^2\|_1$, where 524 $c(\mathbf{u}) = \|\mathbf{u}\|_1 : \mathbb{R}^M \to \mathbb{R}$, and $\mathbf{s}(\mathbf{x}) = \frac{1}{M}(\mathbf{\check{z}} - |\mathbf{H}\mathbf{x}|^2) : \mathbb{C}^N \to 525$ \mathbb{R}^M . It is easy to verify for any $\mathbf{u}, \mathbf{v} \in \mathbb{R}^N$ that it holds 526

$$|c(\mathbf{u}) - c(\mathbf{v})| \le \sum_{n=1}^{M} |u_m - v_m| = \|\mathbf{u} - \mathbf{v}\|_1 \le \sqrt{M} \|\mathbf{u} - \mathbf{v}\|_2$$

where the last inequality arises from the equivalence of norms. 527 By definition, this asserts that c is \sqrt{M} -Lipschitz continuous. 528 Focusing now on the complex Jacobian $\nabla_x \mathbf{s}$ for any \mathbf{x} and 529 $\mathbf{y} \in \mathbb{C}^N$, we deduce 530

$$\left\|\nabla_{\mathbf{x}}\mathbf{s}(\mathbf{x}) - \nabla_{\mathbf{x}}\mathbf{s}(\mathbf{y})\right\|_{2} = \frac{1}{M} \left\|\mathbf{H}^{\mathcal{H}}\mathbf{H}(\mathbf{x}-\mathbf{y})\right\|_{2} \le L \|\mathbf{x}-\mathbf{y}\|_{2}$$

with $L := (2/M)\lambda_{\max}(\mathbf{H}^{\mathcal{H}}\mathbf{H})$, which confirms that $\nabla_x \mathbf{s}$ is *L*-531 Lipschitz continuous. 532

Appealing to the results in [26, Th. 5.3], one can conclude 533 that our deterministic prox-linear SE solver with constant step 534 size $\mu \leq 1/(L\sqrt{M}) = \sqrt{M}/(2\lambda_{\max}(\mathbf{H}^{\mathcal{H}}\mathbf{H}))$ converges to a 535 stationary point of $\ell(\mathbf{x})$ in (34). 536

We provide in the sequel conditions on the function ℓ such 537 that exact recovery of power system states by our prox-linear 538 SE solvers is guaranteed. Going beyond [21], which is limited 539 to optimization over real-valued variables, we introduce two 540 complimentary conditions on $\ell(\mathbf{x})$ and its linearization $\ell_{\mathbf{x}}(\mathbf{y})$ of 541 complex-valued variables $\mathbf{x}, \mathbf{y} \in \mathbb{C}^N$. 542

If ℓ is defined over real-valued vectors, namely $\ell(\mathbf{x}) : \mathbb{R}^N \to$ 543 \mathbb{R} , we can readily have the following stability condition for es-544 tablishing fast convergence: for any given $\mathbf{v} \in \mathbb{R}^N$, there exists 545 a constant $\rho > 0$ such that $\ell(\mathbf{x}) - \ell(\mathbf{v}) \ge \rho \|\mathbf{x} - \mathbf{v}\|_2 \|\mathbf{x} + \mathbf{v}\|_2$ 546 holds for all $\mathbf{x} \in \mathbb{R}^N$, which is included in Appendix B for 547 completeness. This condition, though crucial for establishing 548 fast convergence, does not generalize to functions of complex-549 valued variables. It is obvious in the complex case that if 550 $\mathbf{v} \in \mathbb{C}^N$ is an optimal solution to (34), then $e^{j\phi}\mathbf{v}$ with any $\phi \in$ 551 $[0, 2\pi)$ is also an optimal solution. This ambiguity thus prompts 552 us to define the Euclidean distance of any estimate $\mathbf{x} \in \mathbb{C}^N$ to 553 $\mathbf{v} \in \mathbb{C}^N$ as dist $(\mathbf{x}, \mathbf{v}) := \min_{\phi \in [0, 2\pi)} \|\mathbf{x} - e^{j\phi} \mathbf{v}\|_2$. This def-554 inition can be thought of as enforcing zero-phase angle at the 555 reference bus, a standard procedure in power flow problems 556 to eliminate the phase ambiguity. Invoking a result for stable 557 phase retrieval in [31, Th. 3.1], we define the following stability 558 condition for functions $\ell(\mathbf{x}) : \mathbb{C}^N \to \mathbb{R}$, which also applies to 559 practical settings where the measurements may contain noise or 560 be even adversarially corrupted. 561

562 **Condition 1:** There exists some constant $\rho > 0$ such that 563 the inequality holds for all $\mathbf{x} \in \mathbb{C}^N$

$$\ell(\mathbf{x}) - \ell(\mathbf{v}) \ge \rho \sqrt{\|\mathbf{x} - \mathbf{v}\|_2^2 \|\mathbf{x} + \mathbf{v}\|_2^2 - 4|\Im(\mathbf{x}^{\mathcal{H}}\mathbf{v})|^2}.$$
 (41)

Evidently, when the measurements $\check{\mathbf{z}} = |\mathbf{H}\mathbf{v}|^2$ are noiseless, it holds that $\ell(\mathbf{v}) = 0$. Similar to the real-valued case studied in [21], our Condition 1 is instrumental in establishing fast convergence of the prox-linear algorithm for optimizing functions of complex-valued variables. Besides Condition 1, we require a condition on the linearization $\ell_{\mathbf{x}}(\mathbf{y}) : \mathbb{C}^N \to \mathbb{R}$ of $\ell(\mathbf{x}) = c(\mathbf{s}(\mathbf{x}))$ around \mathbf{x} defined by

$$\ell_{\mathbf{x}}(\mathbf{y}) := c \big(\mathbf{s}(\mathbf{x}) + 2 \Re \big(\nabla_{\mathbf{x}}^{\mathcal{H}} \mathbf{s}(\mathbf{x}) (\mathbf{y} - \mathbf{x}) \big) \big) \,. \tag{42}$$

571 **Condition 2:** There exists a constant $L < +\infty$ such that the 572 inequality holds for all $\mathbf{x}, \mathbf{y} \in \mathbb{C}^N$

$$|\ell(\mathbf{y}) - \ell_{\mathbf{x}}(\mathbf{y})| \le \frac{L}{2} \|\mathbf{x} - \mathbf{y}\|_{2}^{2}.$$
(43)

This condition basically requires that the locally linearized convex approximation $\ell_{\mathbf{x}}(\mathbf{y})$ is quadratically close [cf. (43)] to the nonconvex function $\ell(\mathbf{y})$. Indeed, the ℓ_1 -based PSSE cost function ℓ in (34) automatically satisfies Condition 2 globally. To show this, let us first express $\ell_{\mathbf{x}}(\mathbf{y})$ according to the definition of (42)

$$\ell_{\mathbf{x}}(\mathbf{y}) = \frac{1}{M} \sum_{m=1}^{M} \left| |\mathbf{h}_{m}^{\mathcal{H}} \mathbf{x}|^{2} - \check{z}_{m} + 2\Re \left(\mathbf{x}^{\mathcal{H}} \mathbf{h}_{m} \mathbf{h}_{m}^{\mathcal{H}} (\mathbf{y} - \mathbf{x}) \right) \right|.$$

579 On the other hand, for any $m \in \{1, 2, ..., M\}$ and $\mathbf{y} \in \mathbb{C}^N$, 580 the following holds true

$$\left|\mathbf{h}_{m}^{\mathcal{H}}\mathbf{y}\right|^{2} = \left|\mathbf{h}_{m}^{\mathcal{H}}\mathbf{x}\right|^{2} + 2\Re\left(\mathbf{x}^{\mathcal{H}}\mathbf{h}_{m}\mathbf{h}_{m}^{\mathcal{H}}(\mathbf{y}-\mathbf{x})\right) + \left|\mathbf{h}_{m}^{\mathcal{H}}(\mathbf{y}-\mathbf{x})\right|^{2}.$$

Subtracting \check{z}_m from both sides, summing from m = 1 to M, 581 and leveraging the triangle inequality, we have that 582

$$\ell(\mathbf{y}) = \frac{1}{M} \sum_{m=1}^{M} \left| |\mathbf{h}_{m}^{\mathcal{H}} \mathbf{y}|^{2} - \check{z}_{m} \right| \le \ell_{\mathbf{x}}(\mathbf{y}) + \frac{1}{M} \sum_{m=1}^{M} \left| \mathbf{h}_{m}^{\mathcal{H}}(\mathbf{y} - \mathbf{x}) \right|^{2}$$
$$\ell(\mathbf{y}) = \frac{1}{M} \sum_{m=1}^{M} \left| |\mathbf{h}_{m}^{\mathcal{H}} \mathbf{y}|^{2} - \check{z}_{m} \right| \ge \ell_{\mathbf{x}}(\mathbf{y}) - \frac{1}{M} \sum_{m=1}^{M} \left| \mathbf{h}_{m}^{\mathcal{H}}(\mathbf{y} - \mathbf{x}) \right|^{2}.$$

Rewriting the last terms in matrix-vector form yields

$$|\ell(\mathbf{y}) - \ell_{\mathbf{x}}(\mathbf{y})| \le (\mathbf{y} - \mathbf{x})^{\mathcal{H}} \left(\frac{1}{M} \mathbf{H}^{\mathcal{H}} \mathbf{H}\right) (\mathbf{y} - \mathbf{x}) \le \frac{L}{2} \|\mathbf{y} - \mathbf{x}\|_{2}^{2}$$
(44)

which proves that Condition 2 is satisfied globally by the ℓ_1 -loss 584 in (34). Moreover, one can take $L = \lambda_{\max}(\mathbf{H}^{\mathcal{H}}\mathbf{H}/M)$, namely 585 the largest eigenvalue of matrix $\mathbf{H}^{\mathcal{H}}\mathbf{H}/M$. 586

Under Conditions 1 and 2, we now devote to exact recovery guarantees for the deterministic prox-linear SE solver in 588 Algorithm 1. The following result implies exact recovery of 589 power system states at a quadratic rate by our proposed proxlinear SE solver in Algorithm 1 under suitable conditions. 591

Theorem 1: Let Conditions 1 and 2 hold. Assuming that 592 the quadratic program (37) is solved exactly per iteration of 593 Algorithm 1, the successive prox-linear SE iterates \mathbf{x}_t satisfy 594

$$\frac{\operatorname{dist}(\mathbf{x}_t, \mathbf{v})}{\|\mathbf{v}\|_2} \le \frac{\rho}{L} \left(\frac{L}{\rho} \cdot \frac{\operatorname{dist}(\mathbf{x}_0, \mathbf{v})}{\|\mathbf{v}\|_2}\right)^{2^*}.$$
 (45)

For readability, the proof of Theorem 1 is postponed to 595 Appendix C. Regarding Theorem 1, three observations come 596 in order. 597

Remark 1 (Exact recovery): If the initialization \mathbf{x}_0 of the 598 iterations is accurate enough, meaning that it satisfies the con-599 dition $\operatorname{dist}(\mathbf{x}_0, \mathbf{v}) < (\rho/L) \|\mathbf{v}\|_2$, the prox-linear SE solver re-600 covers exactly the true state vector $\mathbf{v} \in \mathbb{C}^N$. In terms of initial-601 izations, there are several approaches for this desideratum, three 602 of which are discussed next. Since power systems are typically 603 operating close to the flat voltage profile 1, it is reasonable to 604 initialize the algorithm with $\mathbf{x}_0 = \mathbf{1}$. Moreover, as the voltage 605 magnitudes at all buses are assumed available, one can use the 606 voltage magnitude vector $\mathbf{x}_0 = |\mathbf{v}|$ as the initializer. Alterna-607 tively, it is also feasible to initialize with the estimate found by 608 solving the linearized dc power flow equations.

Remark 2 (Quadratic convergence rate): When dist(\mathbf{x}_0 , 610 \mathbf{v}) < $(\rho/L) \|\mathbf{v}\|_2$ holds true, our prox-linear SE algorithm con-611 verges quadratically fast to the globally optimal solution of 612 the nonconvex and nonsmooth optimization problem (34). Ex-613 pressed differently, to obtain a solution \mathbf{x}_t of (at most) ϵ -relative 614 error, namely dist(\mathbf{x}_t , \mathbf{v})/ $\|\mathbf{v}\|_2 \le \epsilon$, we must only run Algo-615 rithm 1 for about $\log_2 \log_2(1/\epsilon)$ iterations, or equivalently, solve 616 $\log_2 \log_2(1/\epsilon)$ convex quadratic programs as in (37). This, in practice, amounts to about 5 ~ 10 such quadratic programs. 618

Remark 3: Under the condition dist $(\mathbf{x}_0, \mathbf{v}) < (\rho/L) \|\mathbf{v}\|_2$, 619 it is worth pointing out that Condition 1 can be replaced by a 620 condition requiring only the function $\ell(\mathbf{x})$ to satisfy the inequality (49) locally for all \mathbf{x} within the neighborhood of \mathbf{v} defined 622 by dist $(\mathbf{x}, \mathbf{v}) < (\rho/L) \|\mathbf{v}\|_2$. 623



Fig. 1. Exact recovery performance of Algorithm 1 for the IEEE 14-bus system (noiseless case).

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VI. NUMERICAL TESTS

In this section, we perform a number of numerical tests to 625 evaluate our approach and compare with the "workhorse" LS-626 based Gauss-Newton method. Several power network bench-627 marks including the IEEE 14-, 118-, and 300-bus systems 628 were simulated, following the MATLAB-based toolbox MAT-629 POWER [32], [33]. The Gauss-Newton iterations were im-630 plemented by using the embedded SE function "doSE.m" 631 in MATPOWER. To carefully isolate the relative perfor-632 mance of the iterative algorithms, rather than initialization 633 634 employed, all simulated schemes were initialized with the flat voltage vector (i.e., the all-one vector) in all reported 635 experiments. 636

637 A. Tests With Zero Noise

The first experiment examines the exact recovery and conver-638 gence performance of Algorithm 1 from noiseless data on the 639 IEEE 14-, 118-, and 300-bus test systems. The actual voltage 640 magnitude (in p.u.) and angle (in radians) of each bus were uni-641 formly distributed over [0.9, 1.1], and over $[-0.1\pi, 0.1\pi]$. The 642 voltage magnitude squares at all buses as well as the active power 643 flows across all lines were measured. The quadratic programs 644 in Step 5 of Algorithm 1 were solved by the standard convex 645 programming solver SeDuMi [34] with a constant step size of 646 $\mu_t = 1,000$. Algorithm 1 terminates either when a maximum 647 number 20 of iterations are simulated, or when the normal-648 ized distance between two consecutive iterates becomes smaller 649 than 10^{-10} , namely dist $(\mathbf{x}_t, \mathbf{x}_{t-1})/\sqrt{N} \le 10^{-10}$. A total of 650 100 Monte Carlo (MC) runs were carried out. Figs. 1-3 plot the 651 normalized estimation errors dist $(\mathbf{x}_t, \mathbf{v})/\sqrt{N}$ for the 100 MC 652 realizations on the simulated three systems, whose correspond-653 ing L values are 0.9980, 3.0201, and 6.3102. Furthermore, Fig. 4 654 depicts the convergence of the normalized estimation error of 655 Algorithm 1 for the 100 runs on the 14-bus system. Evidently, 656 Algorithm 1 achieves exact recovery over the 100 runs, and en-657 658 joys quadratic convergence in this noiseless setting, validating our theoretical findings in Theorem 1. 659



Fig. 2. Exact recovery performance of Algorithm 1 for the IEEE 118-bus system (noiseless case).



Fig. 3. Exact recovery performance of Algorithm 1 for the IEEE 300-bus system (noiseless case).



Fig. 4. Quadratic convergence of Algorithm 1 in 100 runs for the IEEE 14-bus system (noiseless case).

B. Tests With Outlying Measurements

One of the claimed advantages of the ℓ_1 -based loss function 661 in (34) is its robustness to outliers. In this test, we evaluate 662 the robustness of Algorithm 1 to measurements with outliers 663

Fig. 5. Recovery performance for the IEEE 118-bus system with 1 outlying measurement (noiseless case).

Monte Carlo runs

Gauss-Newtor

664 in terms of the exact recovery. Concretely, the IEEE 118-bus system with its default voltage profile was simulated. The volt-665 age magnitudes at all buses along with the active and reactive 666 power flows across all lines were measured. Considering the fact 667 that the nodal voltage magnitudes in transmission networks are 668 maintained close to one, we assume that only the power meters 669 670 can be compromised. A total of 100 MC runs were performed. Per run, one power flow meter was randomly compromised, 671 whose measurement was purposefully manipulated and ampli-672 fied to five times its original value. 673

Both the Gauss-Newton and Algorithm 1 were simulated in 674 this experiment, whose corresponding normalized estimation 675 errors for the 100 MC realizations are presented in Fig. 5. It is 676 self-evident from the plots that the Gauss-Newton method is not 677 robust to outlying measurements, whereas our proposed prox-678 linear scheme in Algorithm 1 can identify and automatically 679 reject the bad data, yielding exact recovery of the true states in 680 most cases even under adversarial attacks. 681

682 C. Tests With Additive Noise and Outliers

The third experiment assesses the robust estimation perfor-683 mance of Algorithm 1 relative to Gauss-Newton, in a setting 684 where both additive noises and outliers are present. The large 685 IEEE 300-bus benchmark system with its default voltage pro-686 file was simulated. All active and reactive power flows as well 687 as all voltage magnitudes were measured. Additive noise was 688 independently generated from normal distributions having zero-689 mean and standard deviations 0.004 and 0.008 for the voltage 690 magnitude and line flow measurements, respectively [18]. In ad-691 dition to additive noise, 5% of the entire measured power flows 692 were corrupted uniformly at random with "outliers" drawn in-693 dependently from a Gaussian distribution with zero-mean and 694 standard deviation 5. The normalized estimation errors obtained 695 by the Gauss-Newton method and Algorithm 1 for 100 MC in-696 dependent realizations are reported in Fig. 6. Evidently, our 697 developed algorithm consistently exhibits more robust estima-698



Fig. 6. Estimation performance for the IEEE 300-bus system under additive noise and 5% outlying measurements.

tion performance than LS-based Gauss–Newton against additive 699 noise and outlying measurements. 700

Robust PSSE is approached by minimizing the ℓ_1 -based loss 702 function from the vantage point of composite optimization. To 703 enable efficient algorithms and exact state recovery, the power 704 quantities were first transformed into rank-one measurements. 705 Building on advances in nonconvex and nonsmooth composite 706 optimization, two algorithms were put forth for minimizing the 707 ℓ_1 -based loss of the transformed rank-one measurements. Our 708 algorithms require no tuning of parameters, except for a step 709 size. We also developed "stability conditions" on the ℓ_1 -based 710 loss function such that exact state recovery and quadratic con-711 vergence are guaranteed by our approach in the noiseless case. 712 Simulated tests using three IEEE benchmark networks under 713 different settings validate our theoretical findings, and show-714 case the efficacy of our approach. 715

APPENDIX

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A. Wirtinger's Calculus

Introducing the complex conjugate coordinates $[\mathbf{x}^T \, \overline{\mathbf{x}}^T]^T \in 718$ \mathbb{C}^{2N} , one can rewrite $\mathbf{s}(\mathbf{x}) = \mathbf{s}(\mathbf{x}, \, \overline{\mathbf{x}}) \in \mathbb{C}^M$. It is obvious now 719 that $\mathbf{s}(\mathbf{x}, \, \overline{\mathbf{x}})$ becomes holomorphic in \mathbf{x} for a fixed $\overline{\mathbf{x}}$, and vice 720 versa. This leads to the partial Wirtinger derivatives [28] 721

$$\frac{\partial s_m}{\partial \mathbf{x}} := \left. \frac{\partial s_m(\mathbf{x}, \overline{\mathbf{x}})}{\partial \mathbf{x}} \right|_{\overline{\mathbf{x}} = \text{constant}} = \left[\frac{\partial s_m}{\partial x_1} \frac{\partial s_m}{\partial x_2} \cdots \frac{\partial s_m}{\partial x_N} \right]$$
$$\frac{\partial s_m}{\partial \overline{\mathbf{x}}} := \left. \frac{\partial s_m(\mathbf{x}, \overline{\mathbf{x}})}{\partial \overline{\mathbf{x}}} \right|_{\mathbf{x} = \text{constant}} = \left[\frac{\partial s_m}{\partial \overline{x}_1} \frac{\partial s_m}{\partial \overline{x}_2} \cdots \frac{\partial s_m}{\partial \overline{x}_N} \right]$$

where the partial derivative with respect to \mathbf{x} ($\overline{\mathbf{x}}$) treats $\overline{\mathbf{x}}$ (\mathbf{x}) as a 722 constant in s_m . The complex gradient of $s_m(\mathbf{x}, \overline{\mathbf{x}})$ with respect 723 to \mathbf{x} or $\overline{\mathbf{x}}$ can be defined by 724

$$\nabla_{\mathbf{x}} s_m := \left(\frac{\partial s_m}{\partial \mathbf{x}}\right)^{\mathcal{H}}, \text{ and } \nabla_{\overline{\mathbf{x}}} s_m := \left(\frac{\partial s_m}{\partial \overline{\mathbf{x}}}\right)^{\mathcal{H}}$$

10⁰

10

10

0 10 20 30 40 50 60 70 80 90 100

Normalized error

yielding the complex gradient of s_m in new coordinate system

$$\nabla_c s_m := \left[\nabla_{\mathbf{x}}^{\mathcal{T}} s_m \, \nabla_{\overline{\mathbf{x}}}^{\mathcal{T}} s_m \right]^{\mathcal{T}} = \left[\frac{\partial s_m}{\partial \mathbf{x}} \, \frac{\partial s_m}{\partial \overline{\mathbf{x}}} \right]^{\mathcal{H}}.$$

726 Upon introducing the complex Jacobian

$$\nabla_c \mathbf{s} := [\nabla_c s_1 \, \nabla_c s_2 \, \cdots \, \nabla_c s_M] \in \mathbb{C}^{2N \times M}$$

we can define for given vectors \mathbf{x} and $\Delta \mathbf{x} \in \mathbb{C}^N$ the following 727 first-order Taylor's expansion: 728

$$\mathbf{s}(\mathbf{x} + \Delta \mathbf{x}) \approx \mathbf{s}(\mathbf{x}) + \nabla_{c}^{\mathcal{H}} \mathbf{s}(\mathbf{x}) \left[\frac{\Delta \mathbf{x}}{\Delta \mathbf{x}} \right]$$
$$= \mathbf{s}(\mathbf{x}) + 2\Re \left(\nabla^{\mathcal{H}} \mathbf{s}_{\mathbf{x}}(\mathbf{x}) \Delta \mathbf{x} \right) \in \mathbb{R}^{M \times N}. \quad (46)$$

B. Supporting Results 729

Proposition 2 ([29, Prop. 3]): Given $\mathbf{a} \in \mathbb{C}^N$ and $b \in \mathbb{R}$, 730 the solution of 731

$$\underset{\mathbf{w}\in\mathbb{C}^{N}}{\operatorname{minimize}}\left|b-\Re\left(\mathbf{a}^{\mathcal{H}}\mathbf{w}\right)\right|+\frac{1}{2\mu}\|\mathbf{w}\|_{2}^{2}$$
(47)

can be obtained as $\hat{\mathbf{w}} := \operatorname{proj}_{\mu}(b/\|\mathbf{a}\|_2^2) \cdot \mathbf{a}$, where $\operatorname{proj}_{\mu}(x)$: 732 $\mathbb{R} \times \mathbb{R}_+ \to \mathbb{R}$ is the projection operator that returns the real 733 number in interval $[-\tau, \tau]$ closest to any given $x \in \mathbb{R}$. 734

Condition 3 [21, Condition 1]: For any given $\mathbf{v} \in \mathbb{R}^N$ 735 there exists a parameter $\rho > 0$ such that function $\ell(\mathbf{x})$ satisfies 736 the following for all $\mathbf{x} \in \mathbb{R}^n$: 737

$$\ell(\mathbf{x}) - \ell(\mathbf{v}) \ge \rho \|\mathbf{x} - \mathbf{v}\|_2 \|\mathbf{x} + \mathbf{v}\|_2.$$
(48)

Concerning the lower bound in (41), we have the next result. 738 **Lemma 1:** For any fixed $\mathbf{v} \in \mathbb{C}^N$, the inequality holds for 739 all $\mathbf{x}, \mathbf{v} \in \mathbb{C}^N$ 740

$$\|\mathbf{x} - \mathbf{v}\|_{2}^{2} \|\mathbf{x} + \mathbf{v}\|_{2}^{2} - 4 \left|\Im(\mathbf{x}^{\mathcal{H}}\mathbf{v})\right|^{2} \ge \|\mathbf{v}\|_{2}^{2} \operatorname{dist}^{2}(\mathbf{x}, \mathbf{v}).$$
(49)

Proof: The left-hand-side term of (49) can be rewritten as 741

$$\begin{aligned} \|\mathbf{x} - \mathbf{v}\|_{2}^{2} \|\mathbf{x} + \mathbf{v}\|_{2}^{2} - 4\Im^{2}(\mathbf{x}^{\mathcal{H}}\mathbf{v}) \\ &= \left(\|\mathbf{x}\|_{2}^{2} + \|\mathbf{v}\|_{2}^{2} - 2\Re(\mathbf{x}^{\mathcal{H}}\mathbf{v})\right) \left(\|\mathbf{x}\|_{2}^{2} + \|\mathbf{v}\|_{2}^{2} + 2\Re(\mathbf{x}^{\mathcal{H}}\mathbf{v})\right) \\ &- 4\Im^{2}(\mathbf{x}^{\mathcal{H}}\mathbf{v}) \\ &= \left(\|\mathbf{x}\|_{2}^{2} + \|\mathbf{v}\|_{2}^{2}\right)^{2} - 4\left[\Re^{2}(\mathbf{x}^{\mathcal{H}}\mathbf{v}) + \Im^{2}(\mathbf{x}^{\mathcal{H}}\mathbf{v})\right] \\ &= \left(\|\mathbf{x}\|_{2}^{2} + \|\mathbf{v}\|_{2}^{2}\right)^{2} - 4|\mathbf{x}^{\mathcal{H}}\mathbf{v}|^{2} \\ &= \left(\|\mathbf{x}\|_{2}^{2} + \|\mathbf{v}\|_{2}^{2} + 2|\mathbf{x}^{\mathcal{H}}\mathbf{v}|\right) \left(\|\mathbf{x}\|_{2}^{2} + \|\mathbf{v}\|_{2}^{2} - 2|\mathbf{x}^{\mathcal{H}}\mathbf{v}|\right) \\ &\geq \|\mathbf{v}\|_{2}^{2} \left(\|\mathbf{x}\|_{2}^{2} + \|\mathbf{v}\|_{2}^{2} - 2|\mathbf{x}^{\mathcal{H}}\mathbf{v}|\right) \\ &= \|\mathbf{v}\|_{2}^{2} \operatorname{dist}^{2}(\mathbf{x}, \mathbf{v}). \end{aligned}$$

742 Taking the square root from both sides of the inequality yields the statement of Lemma 1. 743

C. Proof of Theorem 1 744

The proof is based on that of [21, Th. 1], but we here gen-745 eralize its results to function optimization over complex do-746 mains. Observe that the regularized function $g(\mathbf{x}) := \ell_{\mathbf{x}_t}(\mathbf{x}) + \ell_{\mathbf{x}_t}(\mathbf{x})$ 747

 $\frac{L}{2} \|\mathbf{x} - \mathbf{x}_t\|_2^2$ is L-strongly convex in $\mathbf{x} \in \mathbb{C}^N$, and its min- 748 imum is attained at \mathbf{x}_{t+1} [cf. (37)]. The standard optimal- 749 ity conditions for strongly convex minimization confirms that 750 $g(\mathbf{x}_{t+1}) \leq g(e^{j \angle \mathbf{x}_t^{\mathcal{H}} \mathbf{v}} \mathbf{v}) - \frac{L}{2} \|e^{j \angle \mathbf{x}_t^{\mathcal{H}} \mathbf{v}} \mathbf{v} - \mathbf{x}_{t+1}\|_2^2$; that is, 751

$$\ell_{\mathbf{x}_{t}}(\mathbf{x}_{t+1}) + \frac{L}{2} \|\mathbf{x}_{t+1} - \mathbf{x}_{t}\|_{2}^{2} \leq \ell_{\mathbf{x}_{t}} \left(e^{j \angle \mathbf{x}_{t}^{\mathcal{H}} \mathbf{v}} \mathbf{v} \right) \\ + \frac{L}{2} \left\| e^{j \angle \mathbf{x}_{t}^{\mathcal{H}} \mathbf{v}} \mathbf{v} - \mathbf{x}_{t} \right\|_{2}^{2} - \frac{L}{2} \left\| e^{j \angle \mathbf{x}_{t}^{\mathcal{H}} \mathbf{v}} \mathbf{v} - \mathbf{x}_{t+1} \right\|_{2}^{2}.$$
(50)
ecalling now Condition 2, we have that 752

Recalling now Condition 2, we have that

$$\ell(\mathbf{x}_{t+1}) \le \ell_{\mathbf{x}_t}(\mathbf{x}_{t+1}) + \frac{L}{2} \|\mathbf{x}_{t+1} - \mathbf{x}_t\|_2^2$$
(51a)

$$\ell_{\mathbf{x}_{t}}\left(\mathrm{e}^{j\angle\mathbf{x}_{t}^{\mathcal{H}}\mathbf{v}}\mathbf{v}\right) \leq \ell(\mathrm{e}^{j\angle\mathbf{x}_{t}^{\mathcal{H}}\mathbf{v}}\mathbf{v}) + \frac{L}{2}\left\|\mathrm{e}^{j\angle\mathbf{x}_{t}^{\mathcal{H}}\mathbf{v}}\mathbf{v} - \mathbf{x}_{t}\right\|_{2}^{2}.$$
 (51b)

Substituting (51a) and (51b) into (50) gives rise to

$$\begin{split} \ell(\mathbf{x}_{t+1}) &\leq \ell_{\mathbf{x}_{t}} (\mathrm{e}^{j \angle \mathbf{x}_{t}^{\mathcal{H}} \mathbf{v}} \mathbf{v}) + \frac{L}{2} \left\| \mathrm{e}^{j \angle \mathbf{x}_{t}^{\mathcal{H}} \mathbf{v}} \mathbf{v} - \mathbf{x}_{t} \right\|_{2}^{2} \\ &- \frac{L}{2} \left\| \mathrm{e}^{j \angle \mathbf{x}_{t}^{\mathcal{H}} \mathbf{v}} \mathbf{v} - \mathbf{x}_{t+1} \right\|_{2}^{2} \\ &\leq \ell(\mathrm{e}^{j \angle \mathbf{x}_{t}^{\mathcal{H}} \mathbf{v}} \mathbf{v}) + L \left\| \mathrm{e}^{j \angle \mathbf{x}_{t}^{\mathcal{H}} \mathbf{v}} \mathbf{v} - \mathbf{x}_{t} \right\|_{2}^{2} - \frac{L}{2} \left\| \mathrm{e}^{j \angle \mathbf{x}_{t}^{\mathcal{H}} \mathbf{v}} \mathbf{v} - \mathbf{x}_{t+1} \right\|_{2}^{2} \end{split}$$

which, in conjunction with $\ell(e^{j \angle \mathbf{x}_t^{\mathcal{H}} \mathbf{v}} \mathbf{v}) = \ell(\mathbf{v})$ in (34), yields 754

$$L \left\| \mathbf{e}^{j \angle \mathbf{x}_{t}^{\mathcal{H}} \mathbf{v}} \mathbf{v} - \mathbf{x}_{t} \right\|_{2}^{2} \ge \ell(\mathbf{x}_{t+1}) - \ell(\mathbf{v}) + \frac{L}{2} \left\| \mathbf{e}^{j \angle \mathbf{x}_{t}^{\mathcal{H}} \mathbf{v}} \mathbf{v} - \mathbf{x}_{t+1} \right\|_{2}^{2}.$$
 (52)

Invoking further Condition 1 in (52), we have that

$$L \left\| e^{j \angle \mathbf{x}_{t}^{\mathcal{H}} \mathbf{v}} \mathbf{v} - \mathbf{x}_{t} \right\|_{2}^{2} \ge \frac{L}{2} \left\| e^{j \angle \mathbf{x}_{t}^{\mathcal{H}} \mathbf{v}} \mathbf{v} - \mathbf{x}_{t+1} \right\|_{2}^{2} + \rho \sqrt{\|\mathbf{x}_{t+1} - \mathbf{v}\|_{2}^{2} \|\mathbf{x}_{t+1} + \mathbf{v}\|_{2}^{2} - 4 \left| \Im(\mathbf{x}_{t+1}^{\mathcal{H}} \mathbf{v}) \right|^{2}}$$
(53)

in which the last term can be replaced with its lower bound 756 $\rho \|\mathbf{v}\|_2 \operatorname{dist}(\mathbf{x}_{t+1}, \mathbf{v})$ established in Lemma 1. Upon dropping 757 the nonnegative term $\frac{L}{2} \| e^{j \angle \mathbf{x}_t^{\mathcal{H}} \mathbf{v}} \mathbf{v} - \mathbf{x}_{t+1} \|_2^2$, and recalling the 758 definition of dist(\mathbf{x}_t, \mathbf{v}), we immediately have 759

$$\rho \|\mathbf{v}\|_2 \operatorname{dist}(\mathbf{x}_{t+1}, \mathbf{v}) \le L \operatorname{dist}^2(\mathbf{x}_t, \mathbf{v})$$

dividing both sides of which by $\rho \|\mathbf{v}\|_2^2$ yields

di

$$\frac{\operatorname{st}(\mathbf{x}_{t+1},\mathbf{v})}{\|\mathbf{v}\|_2} \leq \frac{L}{\rho} \cdot \frac{\operatorname{dist}^2(\mathbf{x}_t,\mathbf{v})}{\|\mathbf{v}\|_2^2} = \frac{\rho}{L} \left(\frac{L}{\rho} \cdot \frac{\operatorname{dist}(\mathbf{x}_t,\mathbf{v})}{\|\mathbf{v}\|_2}\right)^2.$$

Applying the above-mentioned inequality successively from 761 the initialization \mathbf{x}_0 for t iterations through \mathbf{x}_t gives rise to (45), 762 concluding the proof. 763

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Gang Wang (M'18) received the B.Eng. degree in electrical engineering and automation from the Beijing Institute of Technology, Beijing, China, in 2011, and the Ph.D. degree in electrical and computer engineering from the University of Minnesota, Minneapolis, MN, USA, in 2018.

He is currently a Postdoctoral Associate with 870 the Department of Electrical and Computer Engineering, University of Minnesota. His research 872 interests focus on the areas of statistical signal 873

processing, stochastic and (non-)convex optimization with applications 874 to autonomous energy grids, and deep learning. 875

Dr. Wang was a recipient of a National Scholarship (2014), a Guo Rui Scholarship (2017), and an Innovation Scholarship (first place in 2017), all from China, as well as a Best Student Paper Award at the 2017 European Signal Processing Conference.

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Hao Zhu (M'12) received the B.Eng. degree in881electrical engineering from Tsinghua University,882Beijing, China in 2006, and the M.Sc. (minor:883Math) and Ph.D. degrees in electrical engineer-884ing from the University of Minnesota, Minneapo-885lis, MN, USA, in 2009 and 2012, respectively.886

She is currently an Assistant Professor in 887 electrical and computer engineering (ECE) with 888 The University of Texas at Austin, Austin, TX, 889 USA. She was a Postdoctoral Researcher, from 890 2012 to 2013, and then an Assistant Profes- 891

sor in ECE from 2014–2017 with the University of Illinois at Urbana-Champaign. Her current research interests include power system monitoring and operations considering its cyber-physical coupling, as well as energy data analytics.

Dr. Zhu was a recipient of the NSF CAREER Award in 2017, the 896 Siebel Energy Institute Seed Grant, and the U.S. AFRL Summer Faculty Fellowship in 2016. She is the faculty advisor and/or coauthor of two best papers at the North American Power Symposium. She is currently a member of the steering committee of the IEEE Smart Grid representing the IEEE Signal Processing Society. 901

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Georgios B. Giannakis (F'97) received the Diploma degree in electrical engineering from the National Technical University of Athens, Athens, Greece, in 1981, the M.Sc. degree in electrical engineering in 1983, the M.Sc. de gree in mathematics, and the Ph.D. degree in electrical engineering both in 1986 from the University of Southern California, Los Angeles, CA, USA.

He was with the University of Virginia, from

1987 to 1998 and since 1999, he has been a Professor with the University of Minnesota, Minneapolis, MN, USA, where 914 915 he holds an Endowed Chair in Wireless Telecommunications, University of Minnesota McKnight Presidential Chair in Electrical and Com-916 917 puter Engineering, and serves as the Director of the Digital Technology Center. His general interests include the areas of communications, net-918 919 working, and statistical signal processing-subjects on which he has 920 authored/coauthored more than 430 journal papers, 730 conference papers, 25 book chapters, two edited books, and two research monographs 921 922 (h-index 134). His current research focuses on learning from big data, 923 wireless cognitive radios, and network science with applications to social, brain, and power networks with renewables. He is the (co-)inventor 924 925 of 30 patents issued.

Dr. Giannakis was the (co-)recipient of nine best paper awards from 926 927 the IEEE SIGNAL PROCESSING (SP) and Communications Societies, including the G. Marconi Prize Paper Award in Wireless Communications. 928 He is also a recipient of the Technical Achievement Awards from the 929 SP Society (2000), from the EURASIP (2005), a Young Faculty Teach-930 ing Award, the G. W. Taylor Award for Distinguished Research from the 931 932 University of Minnesota, and the IEEE Fourier Technical Field Award 933 (inaugural recipient in 2015). He is a Fellow of the EURASIP, and has served the IEEE in a number of posts, including that of a Distinguished 934 935 Lecturer for the IEEE-SP Society. 936



Jian Sun (M'10) was born in Jilin Province, 937 China, in 1978. He received the Bachelor's de-938 gree from the Department of Automation and 939 Electric Engineering, Jilin Institute of Technol-940 ogy, Changchun, China, in 2001, the Master's 941 degree-from the Changchun Institute of Optics, 942 Fine Mechanics and Physics, Chinese Academy 943 of Sciences (CAS), Changchun, China, in 2004, 944 and the Ph.D. degree-from the Institute of Au-945 tomation, CAS, Beijing, China, in 2007. 946 He was a Research Fellow with the Faculty 947

of Advanced Technology, University of Glamorgan, U.K., from April 2008 948 to October 2009. He was a Postdoctoral Research Fellow with the Beijing Institute of Technology, Beijing, China, from December 2007 to May 950 2010. In May 2010, he joined the School of Automation, Beijing Institute of Technology, where he has been a Professor since 2013. His research 952 interests include control, security, and learning of networked control systems and cyber-physical systems. 954

Dr. Sun is currently on the Editorial Boards of *Journal of Systems* 955 Science and Complexity and Acta Automatica Sinica. 956